

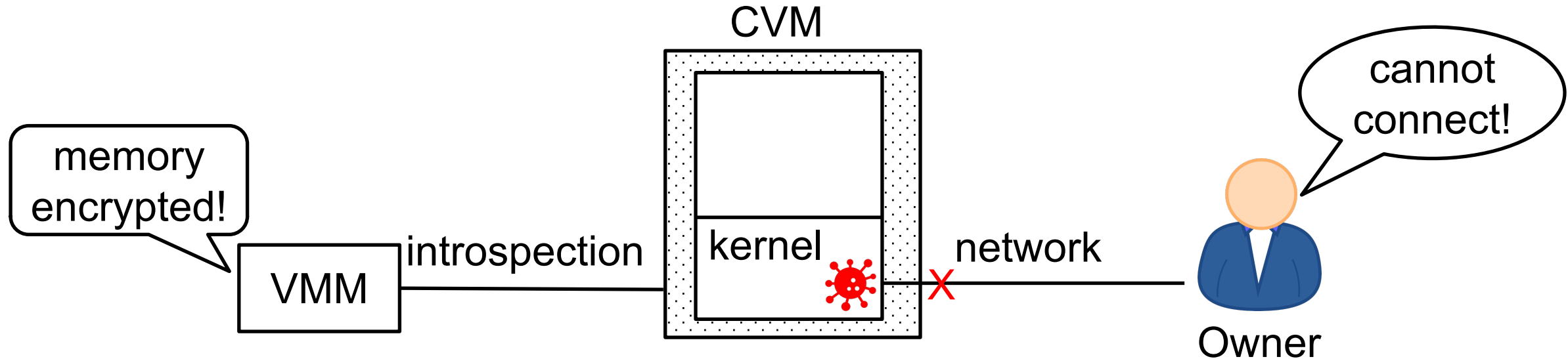
TETD: Trusted Execution in Trust Domains

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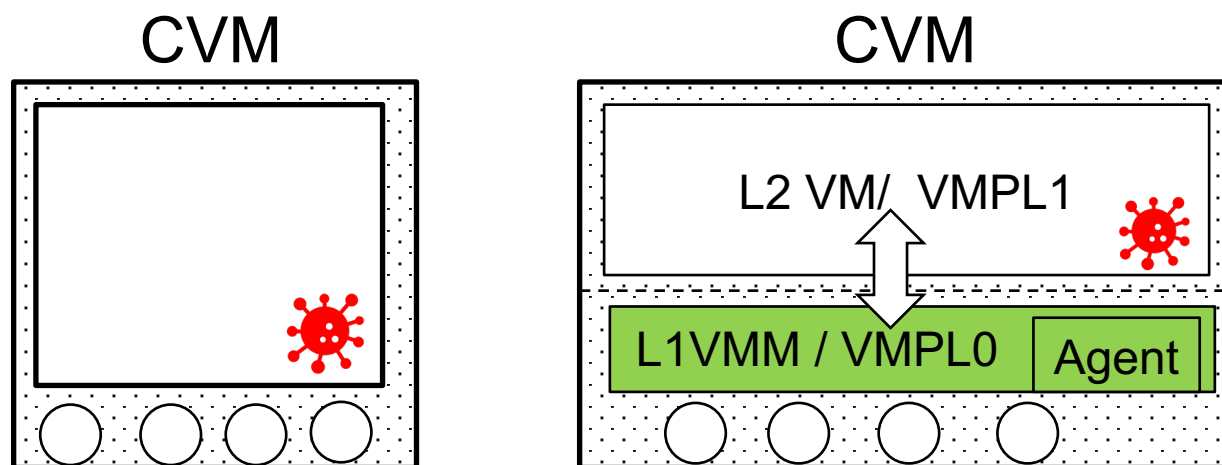
The Problem

- In the event of CVM kernel compromise or crash, the owner loses the foothold to take care of her CVM.



Existing Approach

- **In-VM Privilege Layering**: to insert a **more privileged** and **trusted** layer under the CVM kernel.
 - used in HARDLOG[SP22], Hecate[CCS22], Veil[ASPLOS23], SVSM-vTPM[ACSAC23], 00SEVEN [Sec24], NestedSGX [NDSS25] etc.

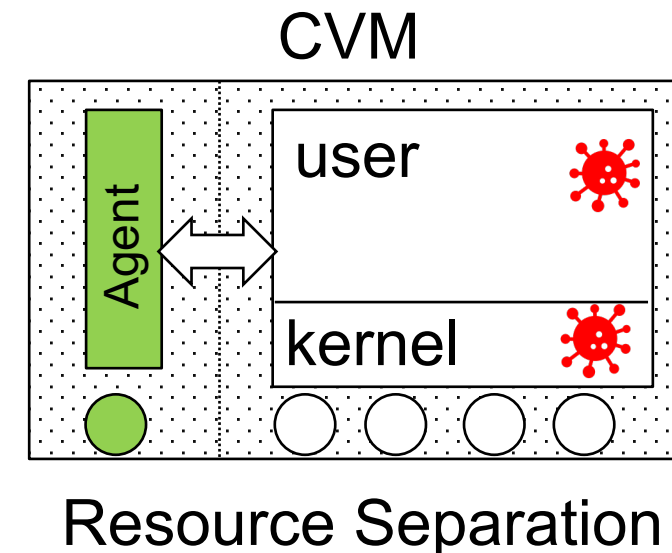
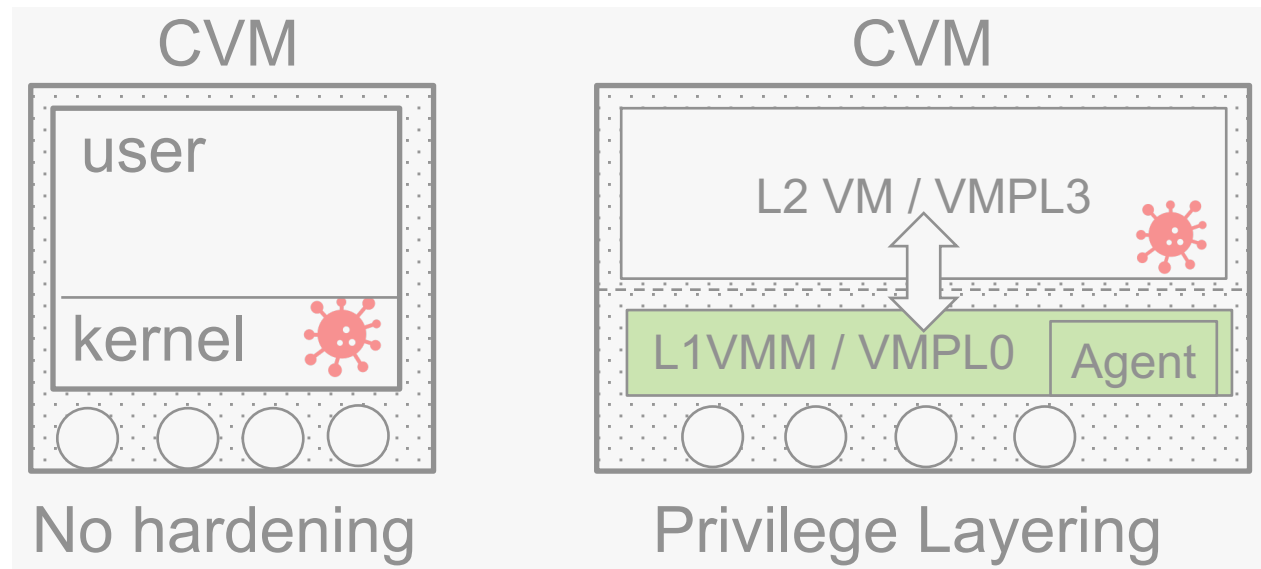


Privilege Layering

Our Approach

- **Resource Separation:**

- A CVM is split into a two (or more) sub-systems with **separated physical memory and vCPUs**.
- To protect a subsystem against another, the untrusted VMM **withholds** the former's resources when others are running.



Results

- **TETD**: a resource-separation based TD hardening scheme without modifying the Intel TDX Module.
- Two execution modes:
 - **exclusive mode** for system-level maintenance.
 - Example case: introspection
 - **collaborative mode** for secure execution against untrusted kernel.
 - Example case: kernel log hardening, enclave-like decryption.
- Pros and Cons:
 - + Secure against full-CVM compromise, including L1VMM or VMPL0 kernel.
 - + No architectural change to CVM; easy to deploy; highly flexible
 - Trust the VMM to faithfully execute the scheme.